

Numerical Optimization

Exercise V

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The Simplex Method

Linear program in standard form

Minimize $f : \mathbb{R}^n \mapsto \mathbb{R}$,

$$f(\mathbf{x}) = \mathbf{c}^\top \mathbf{x},$$

subject to $\mathbf{Ax} = \mathbf{b}$ and $\mathbf{x} \geq 0$, where

$$\mathbf{A} \in \mathbb{R}^{m \times n}, \quad \mathbf{b} \in \mathbb{R}^m, \quad \mathbf{c} \in \mathbb{R}^n, \quad m < n,$$

and \mathbf{A} has full row rank.

Definition (Simplex)

The set

$$\left\{ \mathbf{x} \in \mathbb{R}^n \mid \mathbf{x} = \sum_{i=1}^d \lambda_i \mathbf{x}_i, \sum_{i=1}^d \lambda_i = 1, \lambda_i \geq 0, \mathbf{x}_i \in \mathbb{R}^n \right\}$$

is called d -dimensional *simplex*, $d \in \{0, \dots, n-1\}$, or simply d -*simplex*.

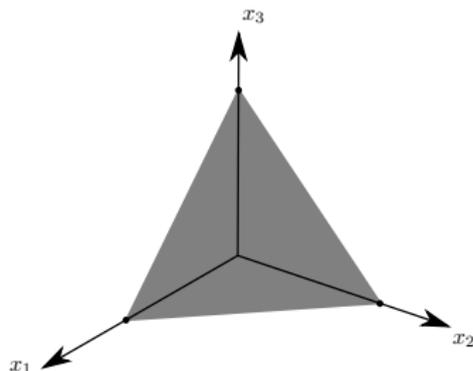
The Simplex Method

So why the name “simplex method”?

Rows of \mathbf{A} define hyperplanes in \mathbb{R}^n . Their intersection restricted to the positive orthant, i.e. the feasible region, is an $n - m$ -simplex (if non-empty and bounded).

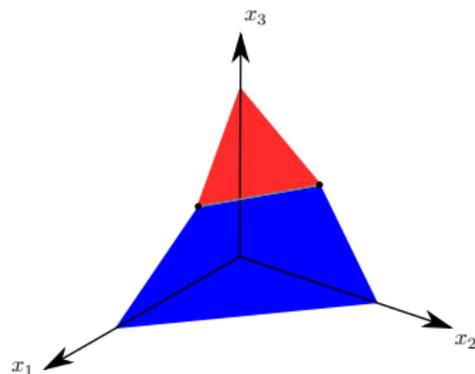
Example: $n = 3$

- ▶ $m = 1$: Picture the infeasible and unbounded case!

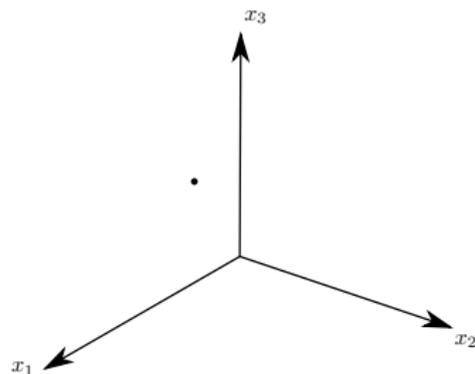


The Simplex Method

- ▶ $m = 2$: If both planes are not entirely infeasible, intersection non-empty due to rank assumption.



- ▶ $m = 3$: At most one feasible point if \mathbf{A} has full row rank \Rightarrow problem trivial.



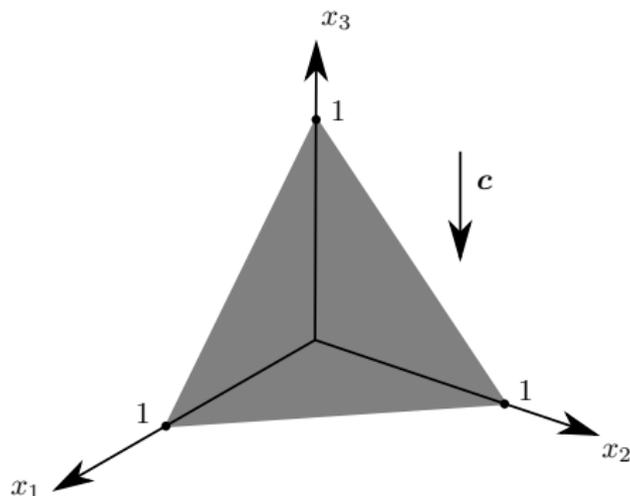
The Simplex Method

Standard simplex gives standard example:

$$\mathbf{A} = \begin{pmatrix} 1 & 1 & 1 \end{pmatrix},$$

$$\mathbf{b} = 1,$$

$$\mathbf{c} = \begin{pmatrix} 0 \\ 0 \\ -1 \end{pmatrix}$$



The Simplex Method

A step-by-step recipe:

1. Start on face of positive orthant, say in the x_1x_3 -plane:

$$\mathcal{B} = \{1\}, \quad \mathcal{N} = \{2, 3\},$$

$$\mathbf{B} = \mathbf{1}, \quad \mathbf{N} = \begin{pmatrix} 1 & 1 \end{pmatrix},$$

$$\mathbf{x}_B = x_1 = \mathbf{B}^{-1}\mathbf{b} = 1, \quad \mathbf{x}_N = \begin{pmatrix} x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}.$$

The Simplex Method

2. Dual variables λ, \mathbf{s} from

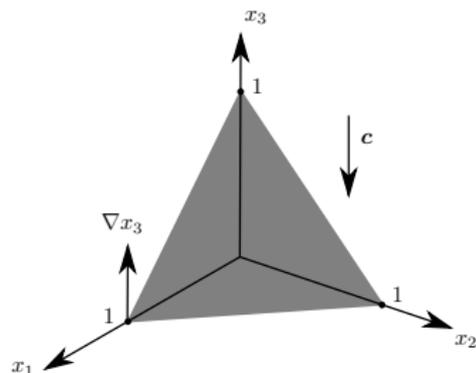
- ▶ complementary slackness $\mathbf{s}^\top \mathbf{x} = 0$

$$\mathbf{s}_B = \mathbf{s}_1 = 0$$

- ▶ Lagrange condition $\mathbf{A}^\top \lambda + \mathbf{s} = \mathbf{c}$:

$$\lambda = \mathbf{B}^{-\top} \mathbf{c}_B = c_1 = 0,$$

$$\mathbf{s}_N = \mathbf{c}_N - \mathbf{N}^\top \lambda = \begin{pmatrix} 0 \\ -1 \end{pmatrix}$$

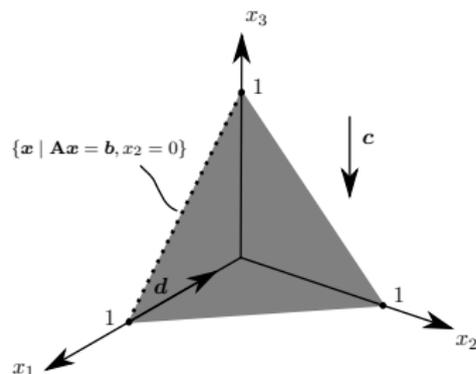


The Simplex Method

3. If $\mathbf{s} \geq \mathbf{0}$, terminate, else select entering index, here $q = 3$.
4. Deactivate $x_q \geq 0$: Increase x_q in \mathbf{x}_N by “decreasing” \mathbf{x}_B in the direction

$$\mathbf{d} = \mathbf{B}^{-1} \mathbf{A}_q = \mathbf{1}$$

(coupling due to $\mathbf{Ax} = \mathbf{b}$).



The Simplex Method

5. Let $j \in \{1, \dots, |\mathcal{B}|\}$ be the row index of \mathbf{x}_B , for which 0 is reached first, i.e. $x_{B,j} - \alpha d_j = 0$ for some $\alpha \in \mathbb{R}$. Turns out that

$$\alpha = \min_{i \in \{1, \dots, |\mathcal{B}|\}} \frac{x_{B,i}}{d_i} = 1 = x_q$$

so that

$$\mathbf{x}_B^+ = 1 - 1 = 0, \quad \mathbf{x}_N^+ = \begin{pmatrix} 0 \\ 1 \end{pmatrix}.$$

6. Map j to leaving index p , update

$$\mathcal{B} = \{p\} = \{3\}, \quad \mathcal{N} = \{q, 2\} = \{1, 2\},$$

$$\mathbf{B} = 1, \quad \mathbf{N} = \begin{pmatrix} 1 & 1 \end{pmatrix},$$

$$\mathbf{x}_B = x_q = x_3 = 1, \quad \mathbf{x}_N = \begin{pmatrix} x_2 \\ x_p \end{pmatrix} = \begin{pmatrix} x_2 \\ x_1 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

(switch x_p in \mathbf{x}_B^+ and x_q in \mathbf{x}_N^+), and continue with 2.

Exercises

1. Consider the linear program $\min -5x_1 - x_2$ subject to

$$\begin{aligned}x_1 + x_2 &\leq 5, \\2x_1 + \frac{1}{2}x_2 &\leq 8, \\x &\geq \mathbf{0}.\end{aligned}$$

- (i) Bring it to standard form.
- (ii) Use the simplex method to solve it.